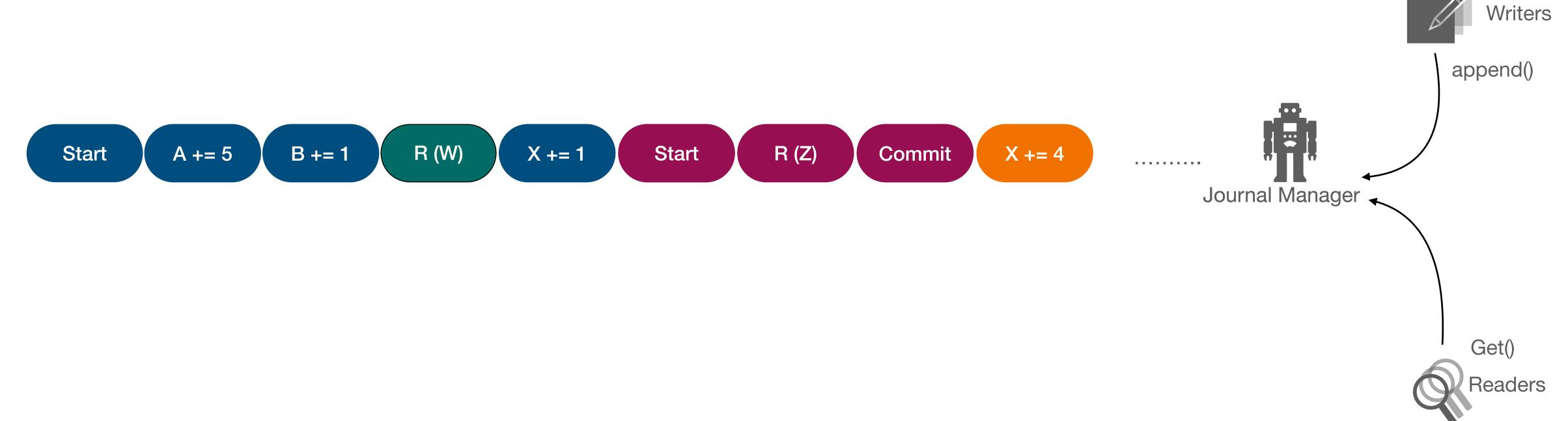
Persisting the Cache in journal based Data-stores Managing Journal Size through truncation

Ayush Pandey
Dr. Annette Bieniusa
Dr. Marc Shapiro

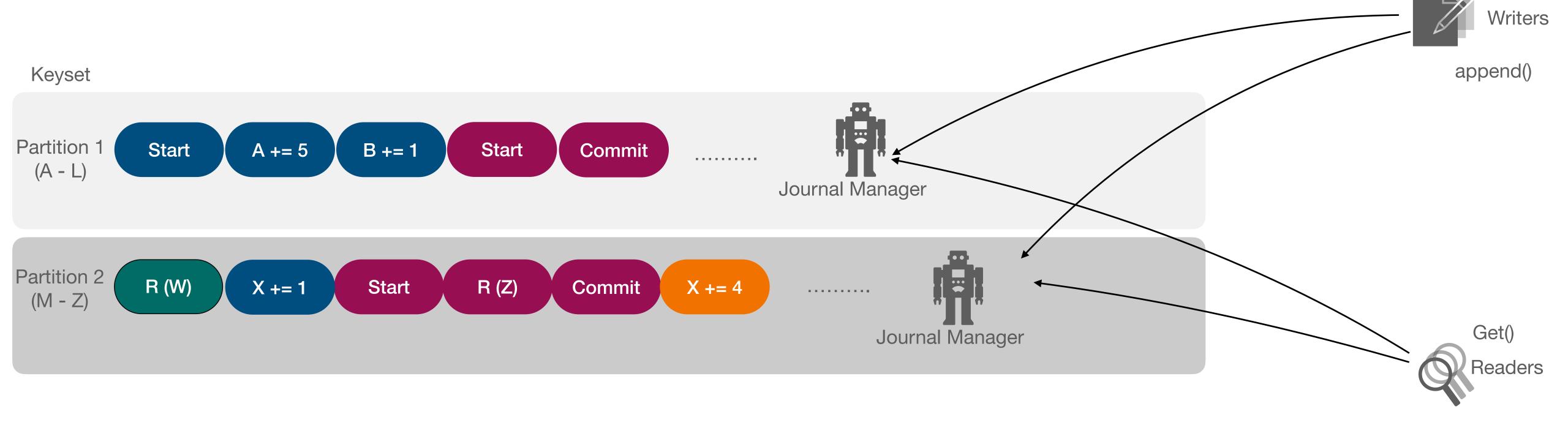
What is a journal?

- Sequence of operations
- Written in a specific order by transactions.
- To be read by transactions to materialise objects.
- The journal can be partitioned to allow distribution.

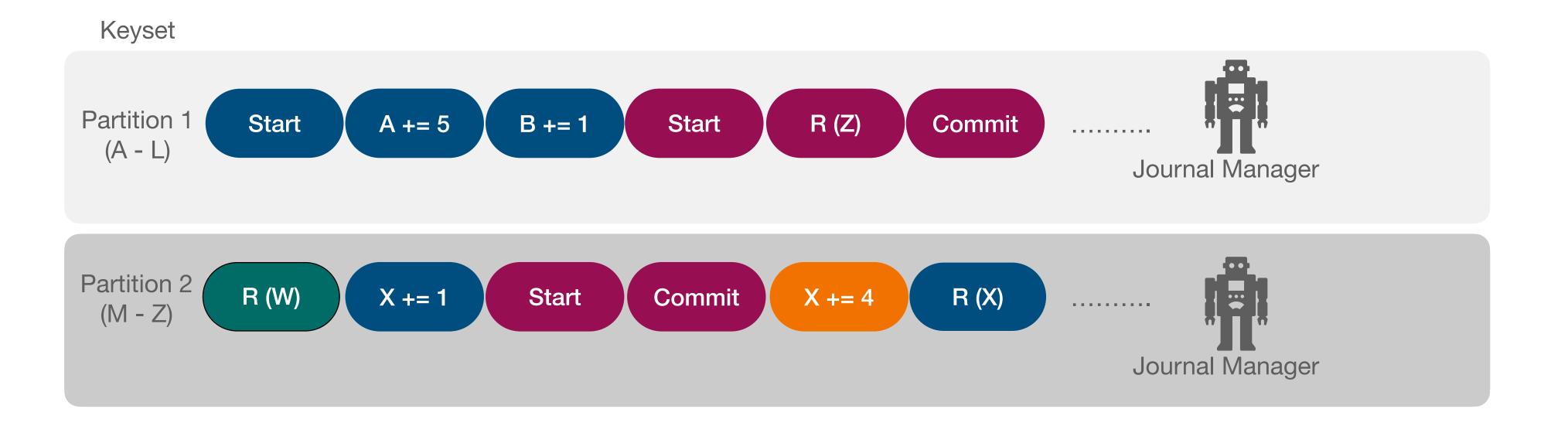


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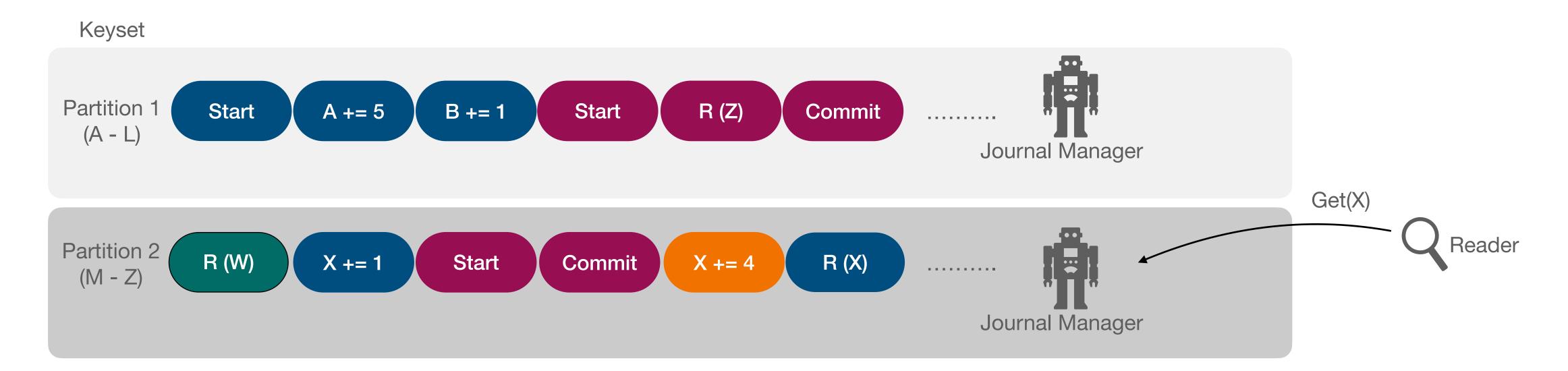


- Read queries lead to materialisations.
- To answer a specific read, the request is sent to the responsible partition.
- The journal manager filters the operations for that key.
- The materialiser uses these filtered operations and creates the object.

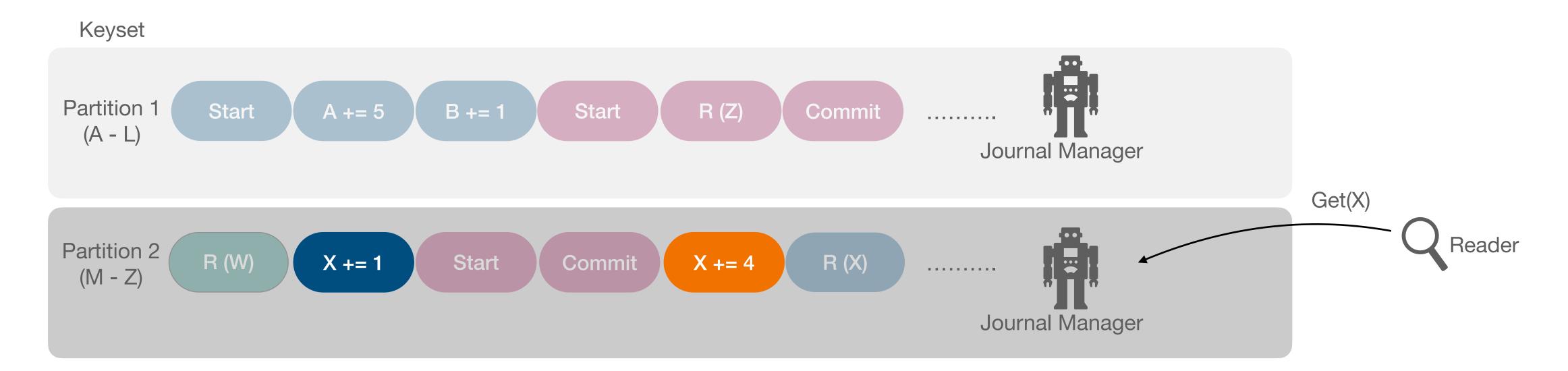




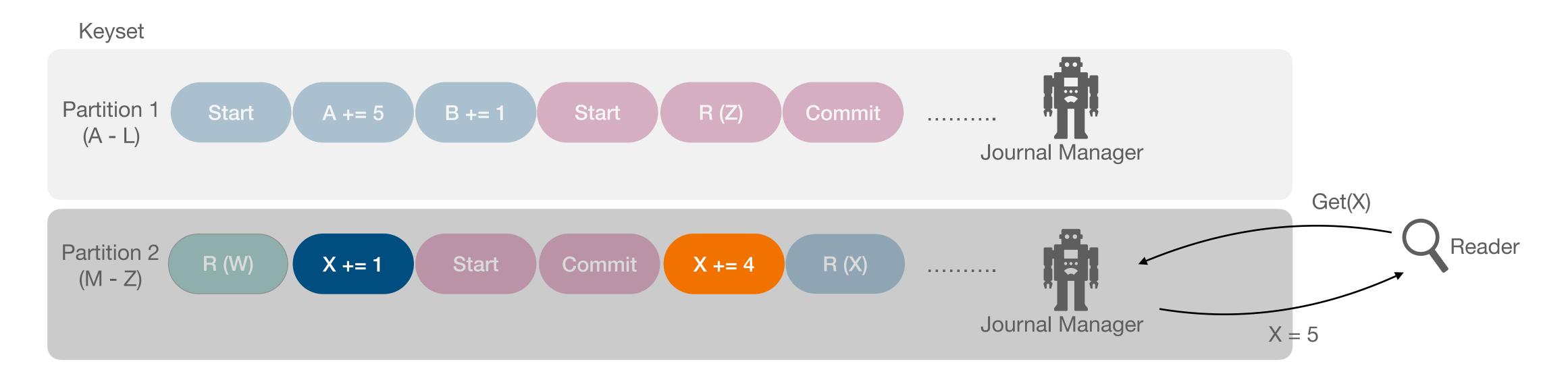
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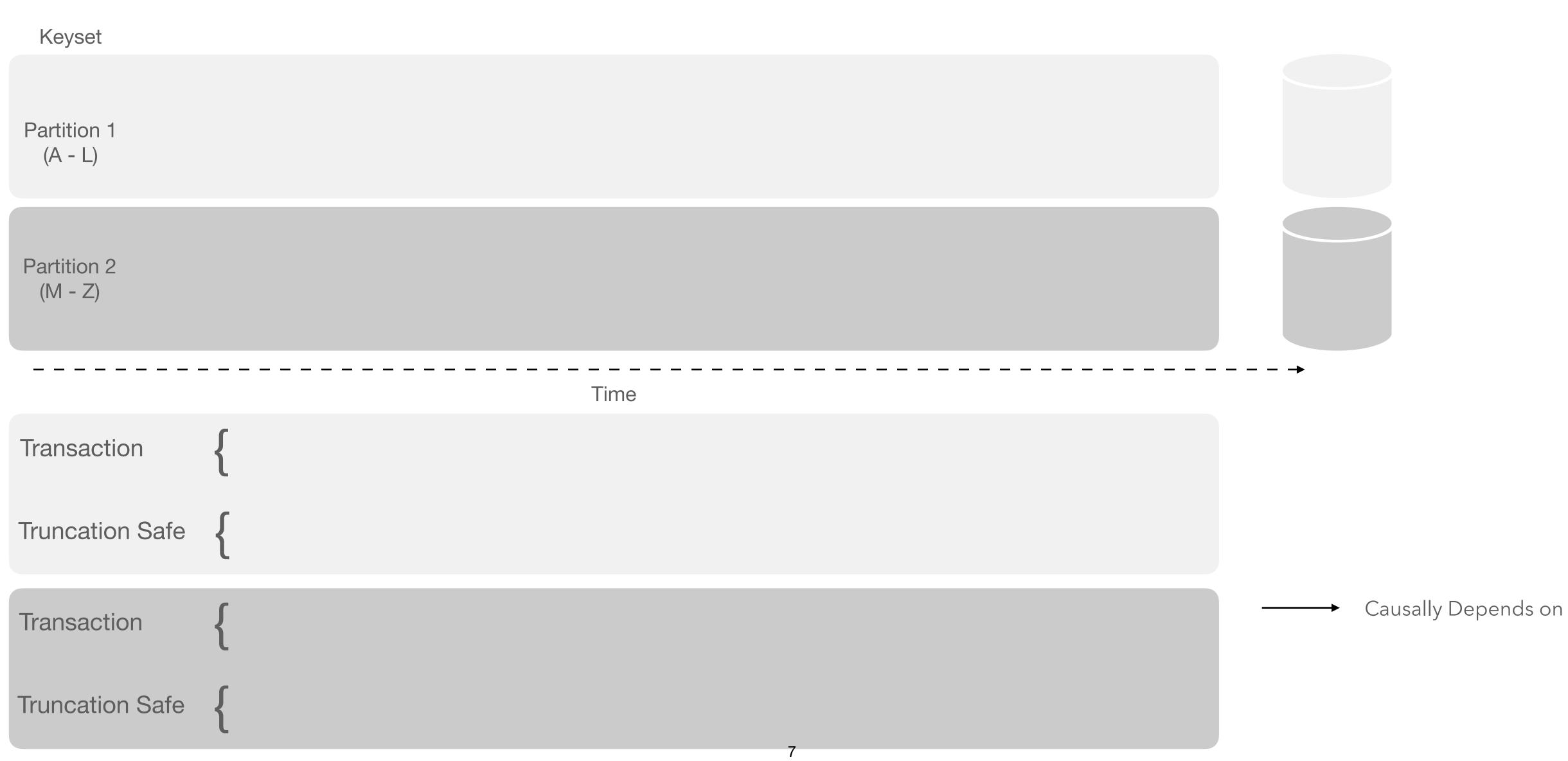
The problem with journals and filtering.

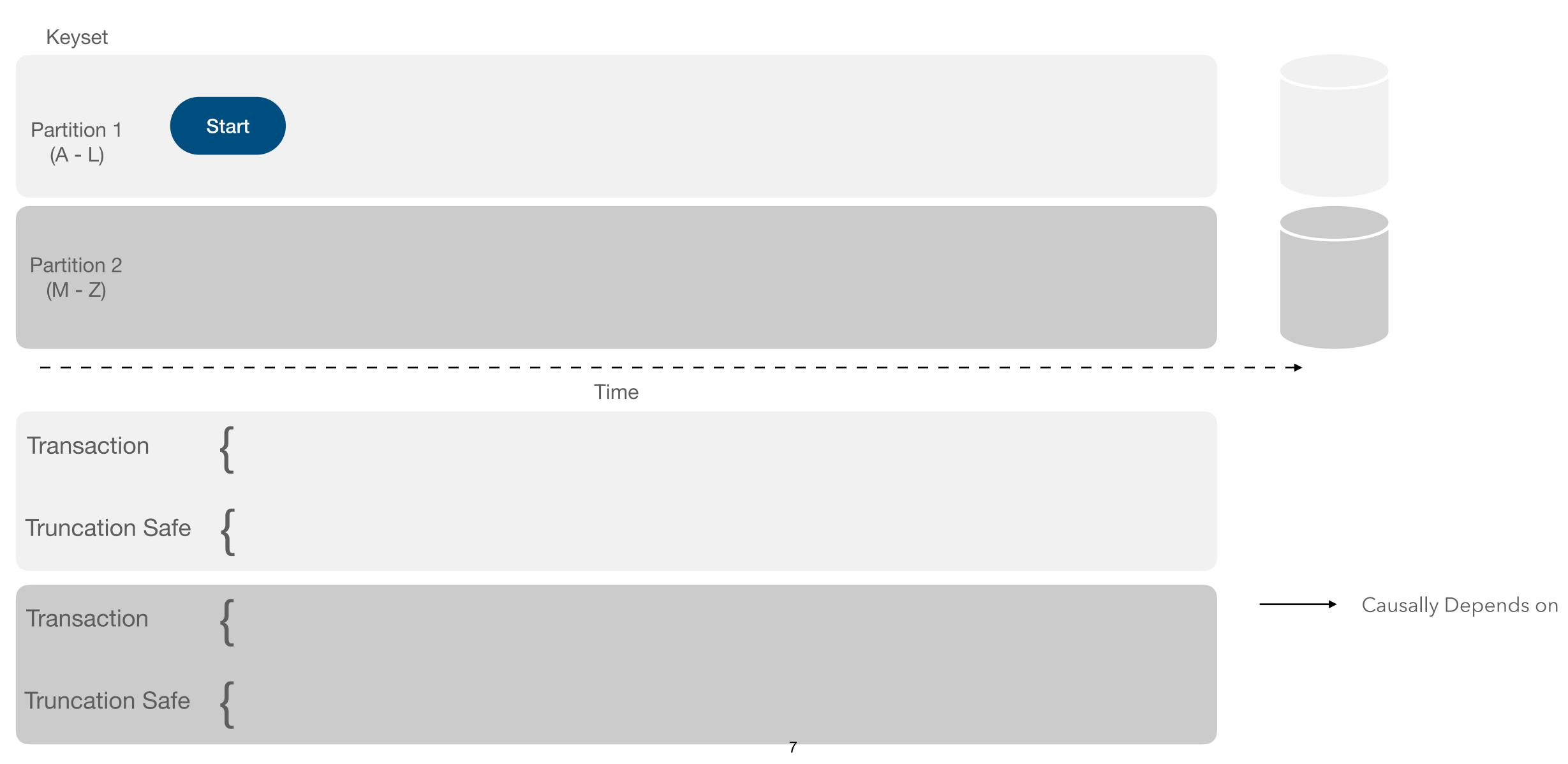
- Each transaction adds and reads some operations to and from the journal.
- The journal is, in the original AntidoteDB design, unbounded and contains all the operations.¹
- We need to filter this large set of read operations to materialise the objects.
- The time required to filter is proportional to the size of the journal.²
- To reduce the time required to filter, reduce the size of the read set.

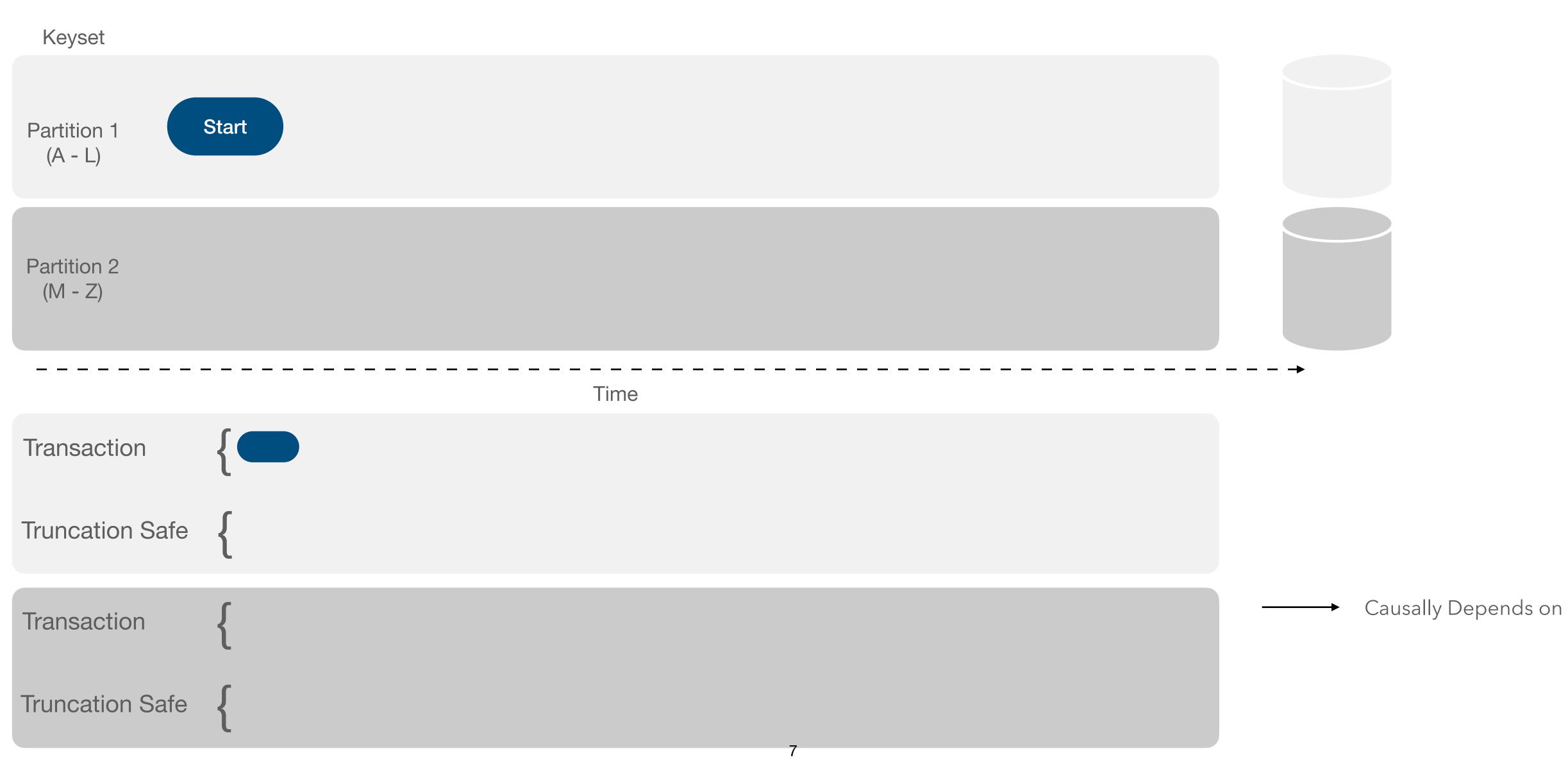
Truncate the journal!

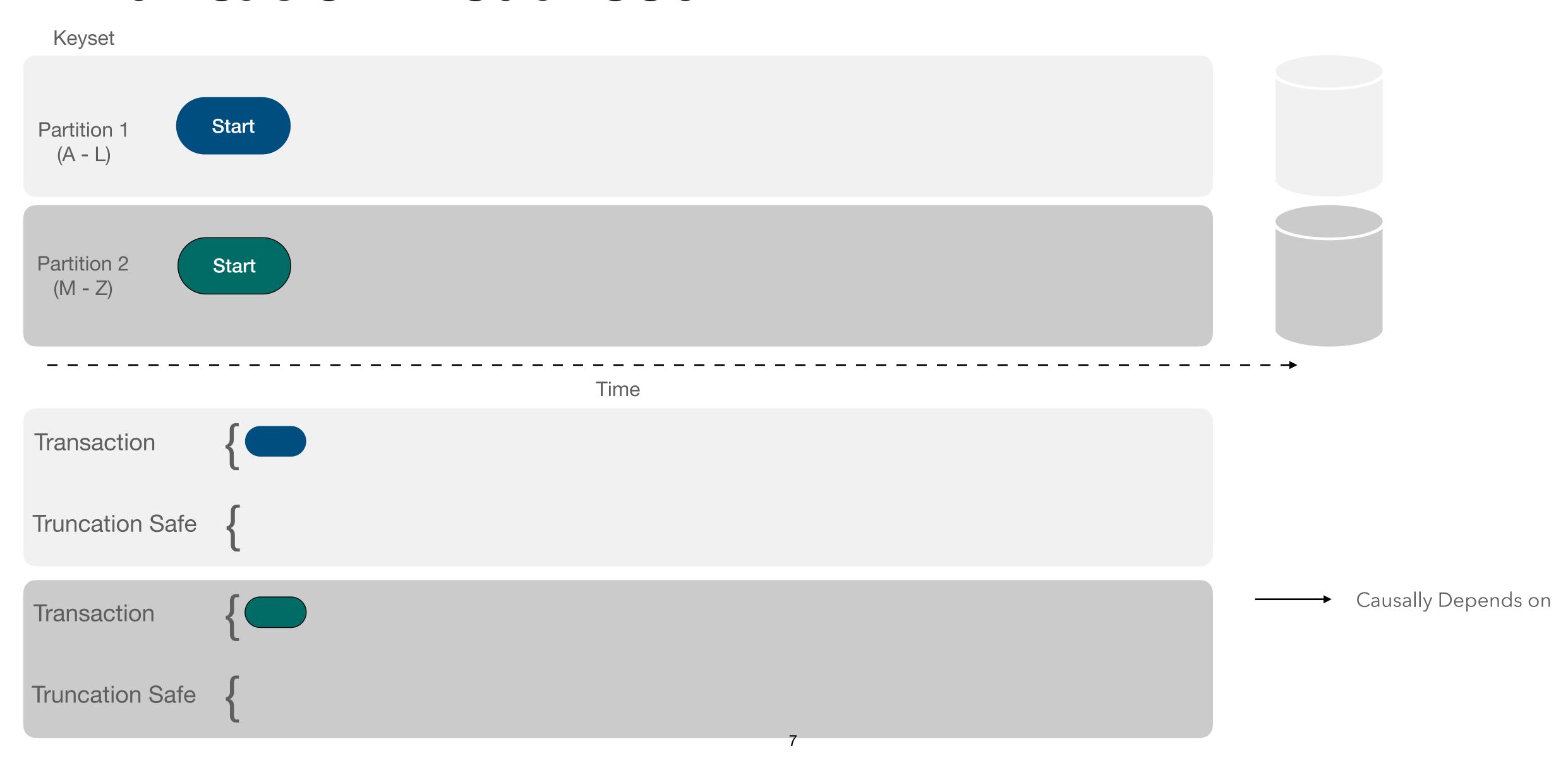
How to truncate a journal?

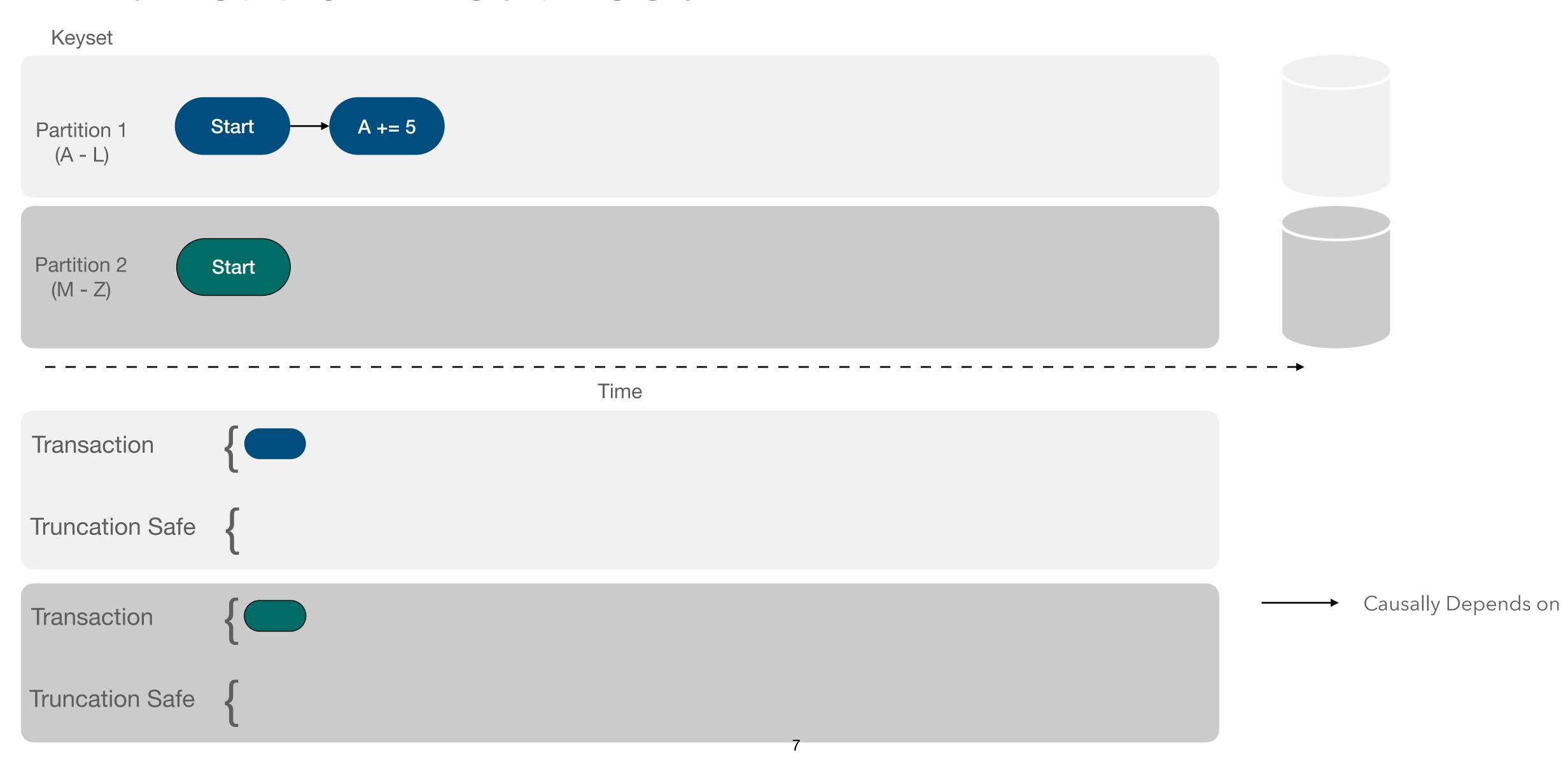
- Journal truncation is destructive.
- We need to maintain causal dependencies between transaction snapshots.
- To ensure safe truncation, there are invariants that need to be maintained.
- Conditions on times of interest during the operation of the database.
- Truncated operations need to be persisted (Check-pointed).

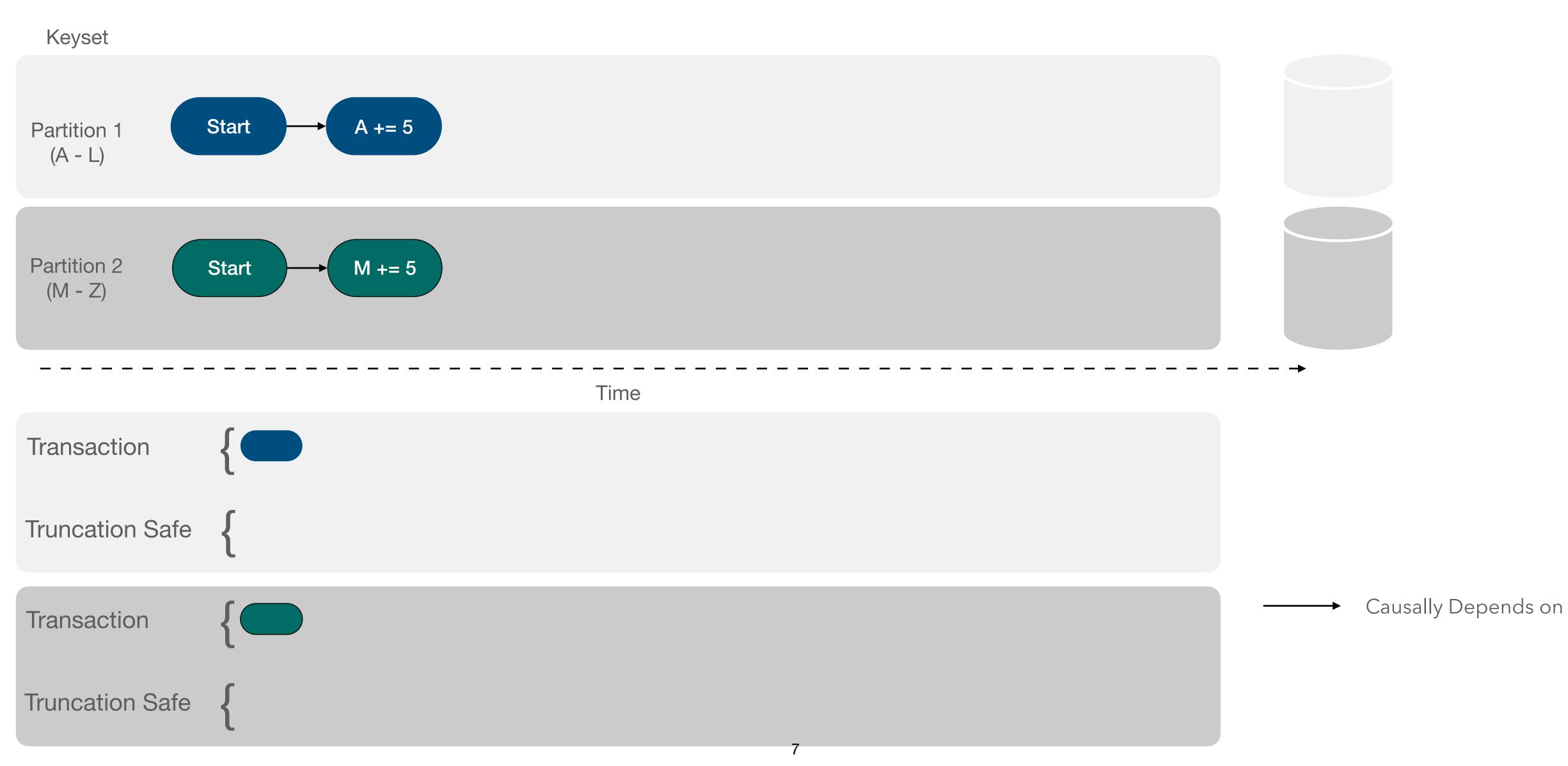


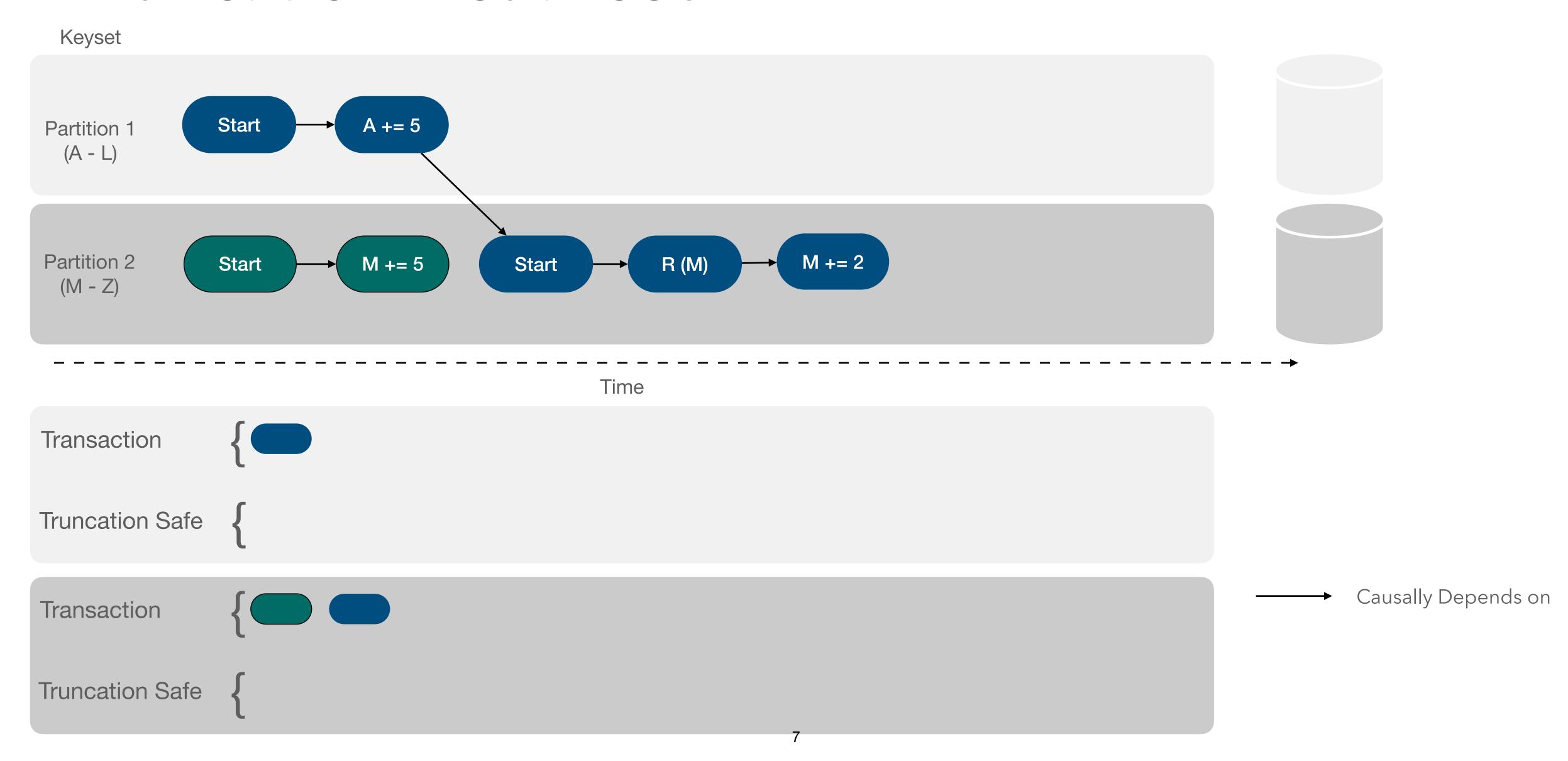


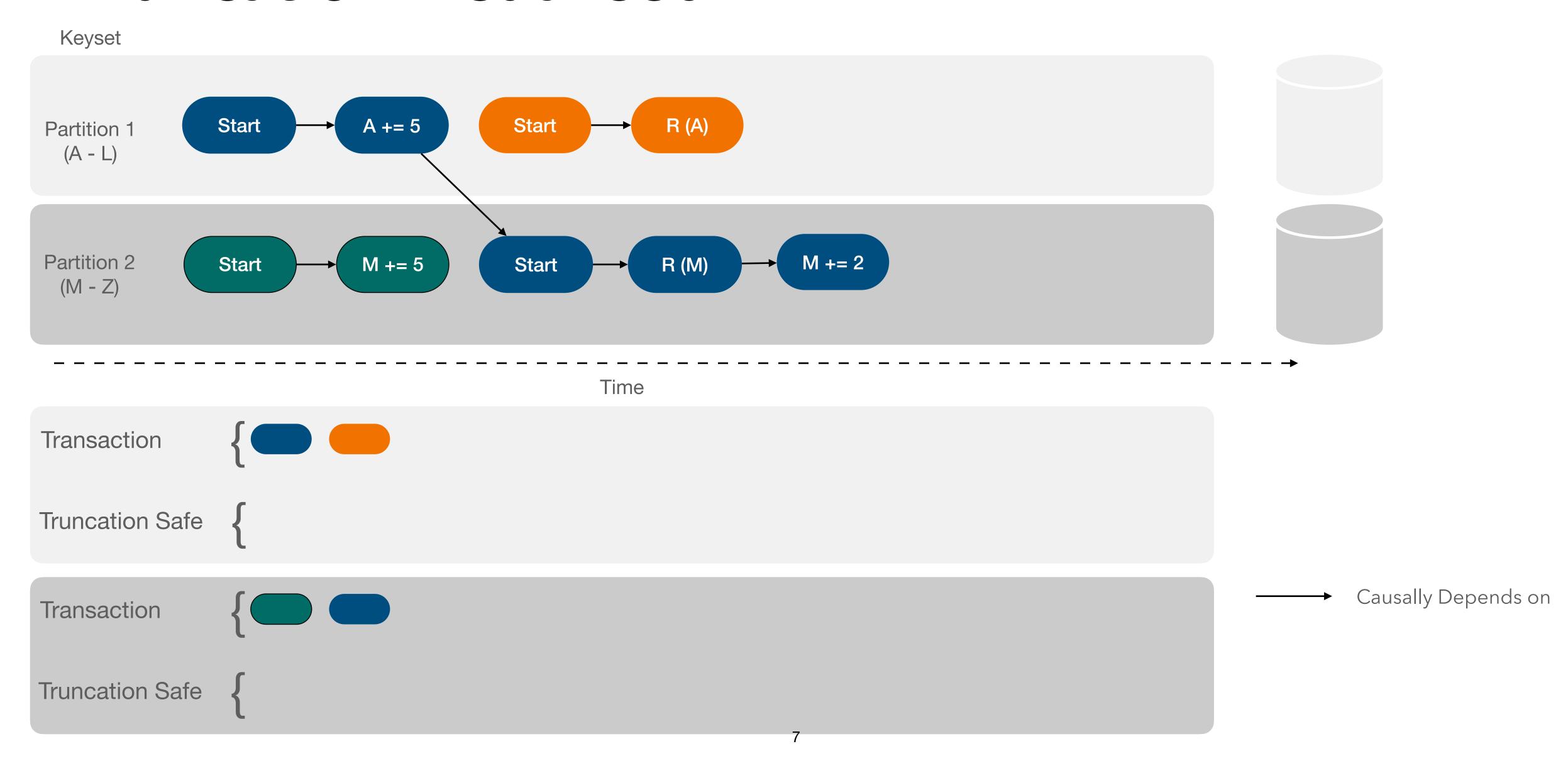


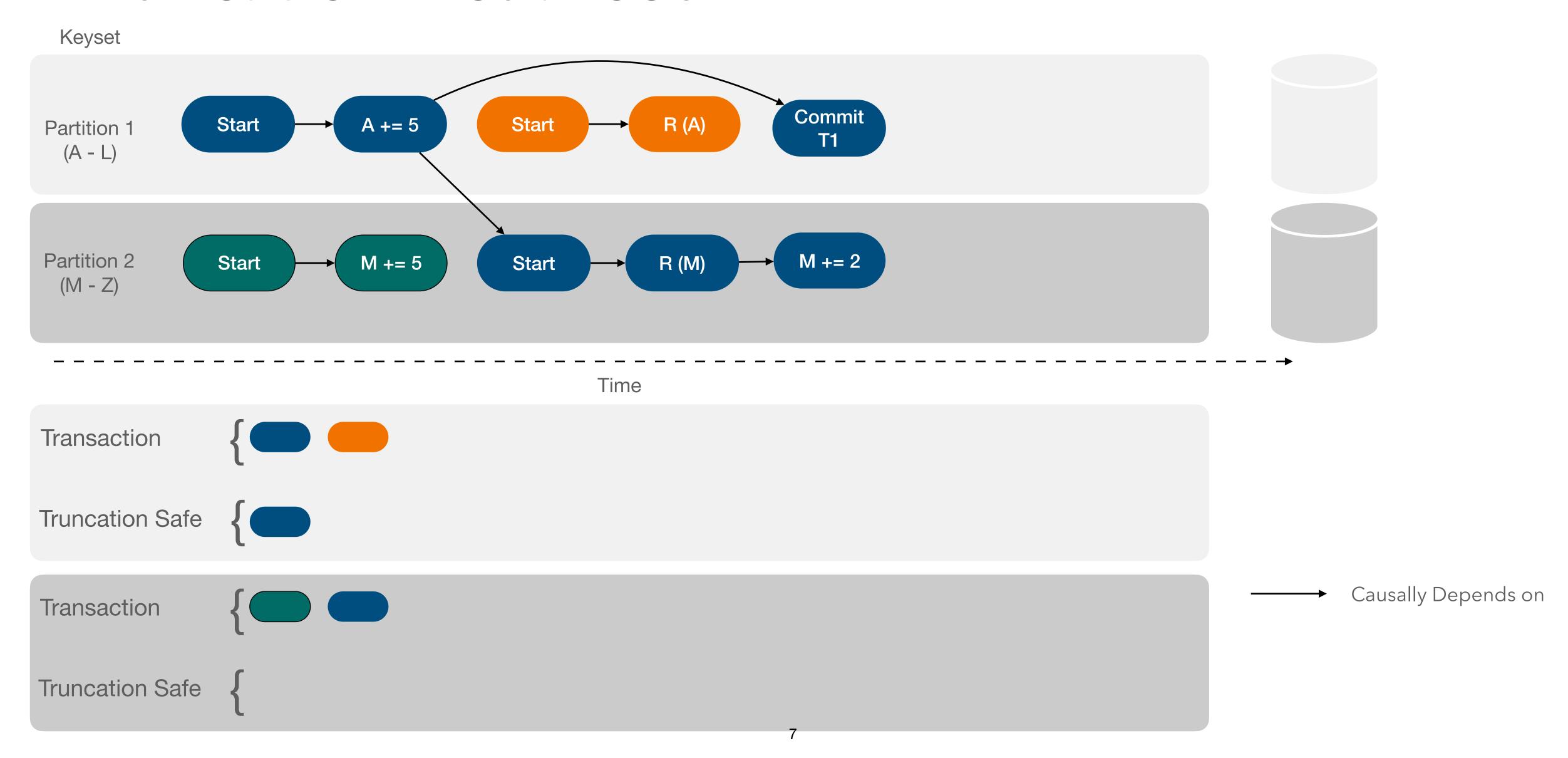


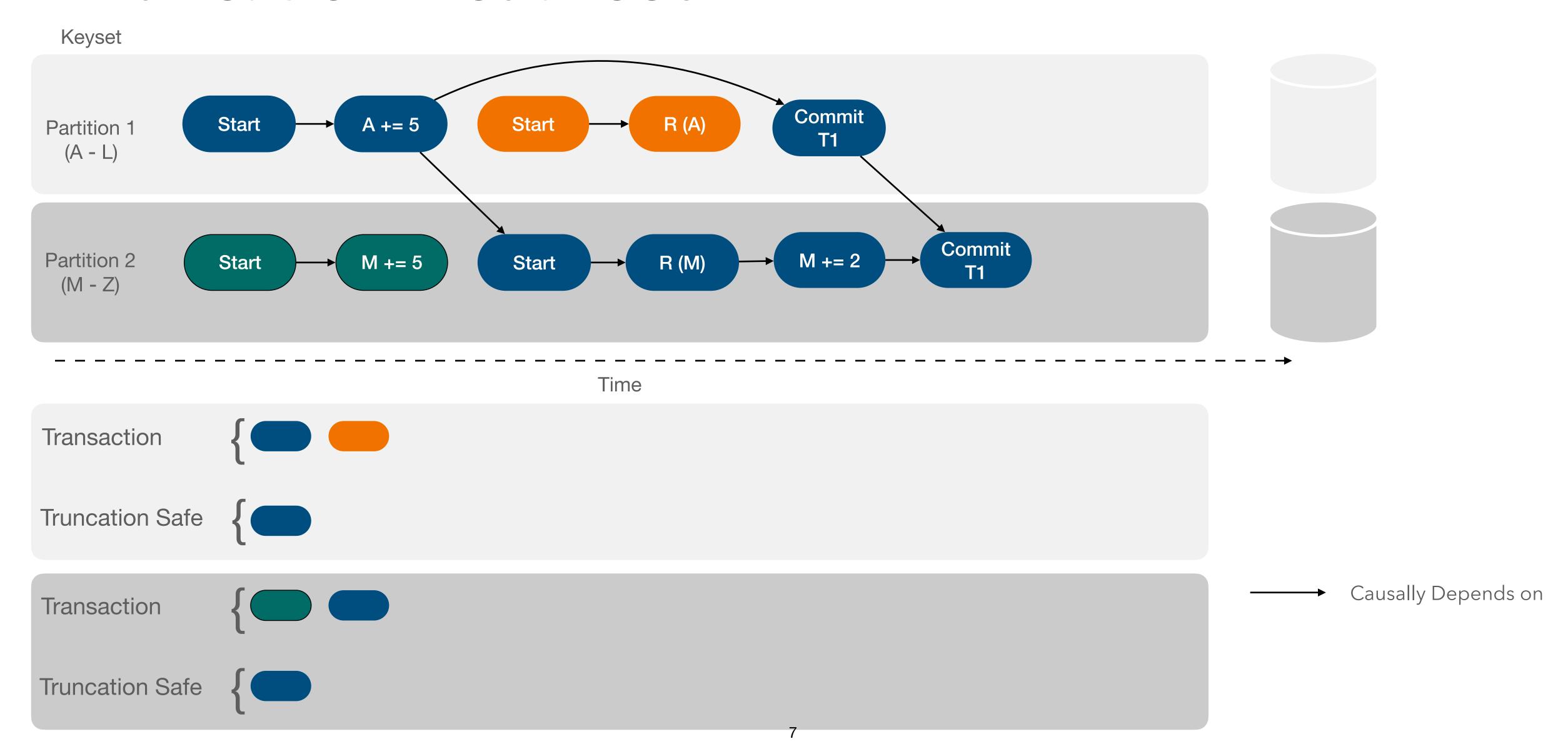


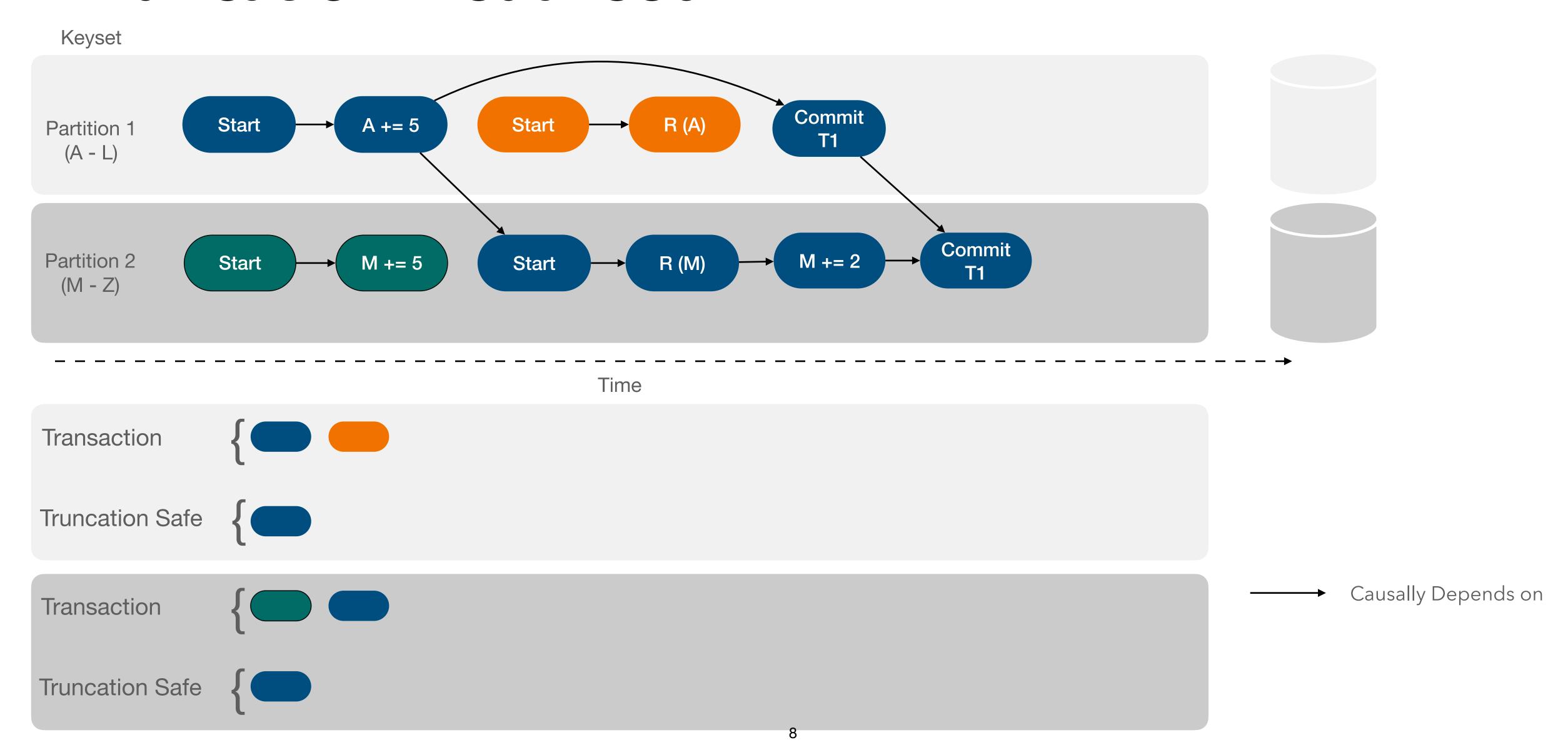


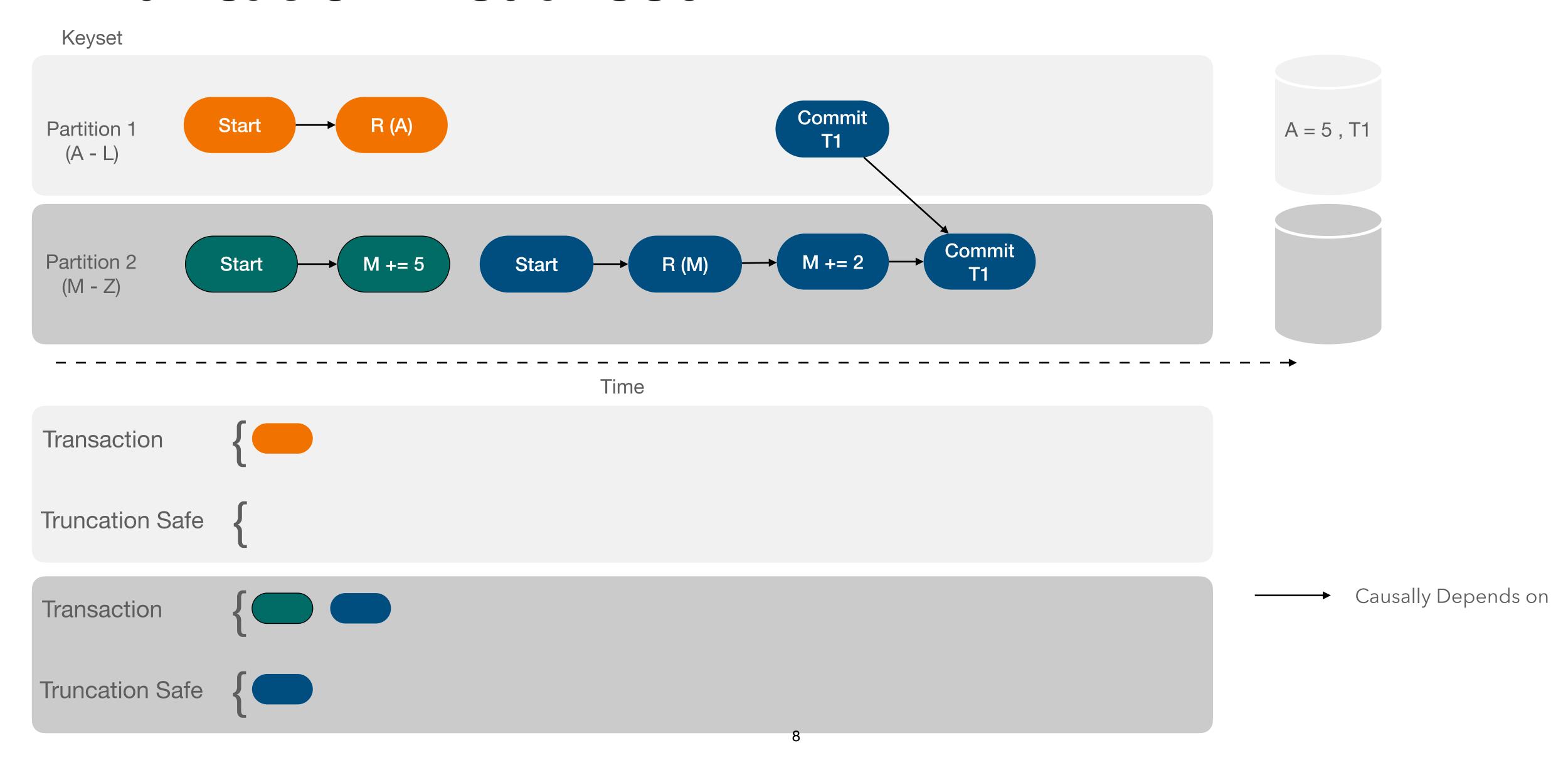


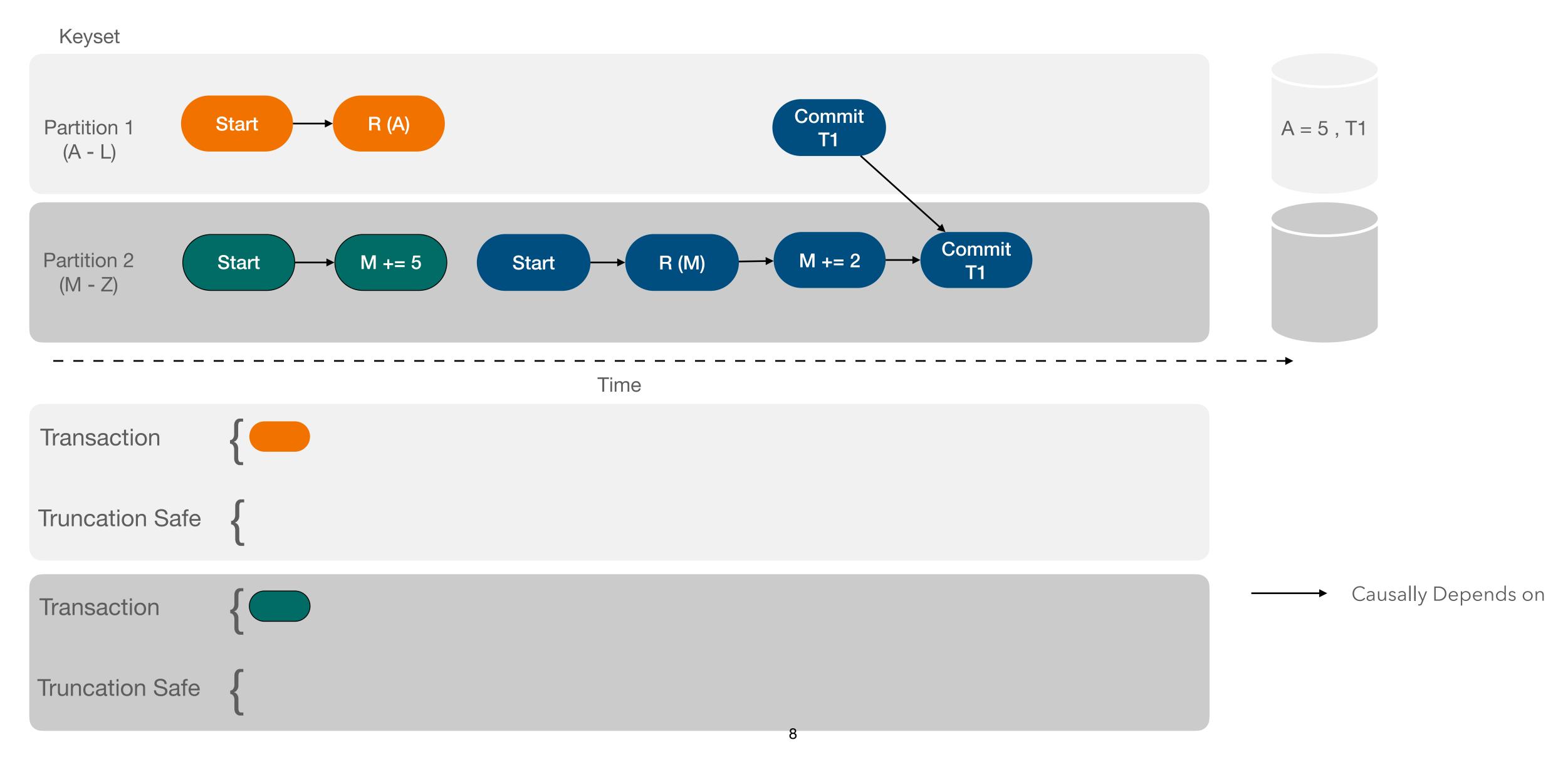


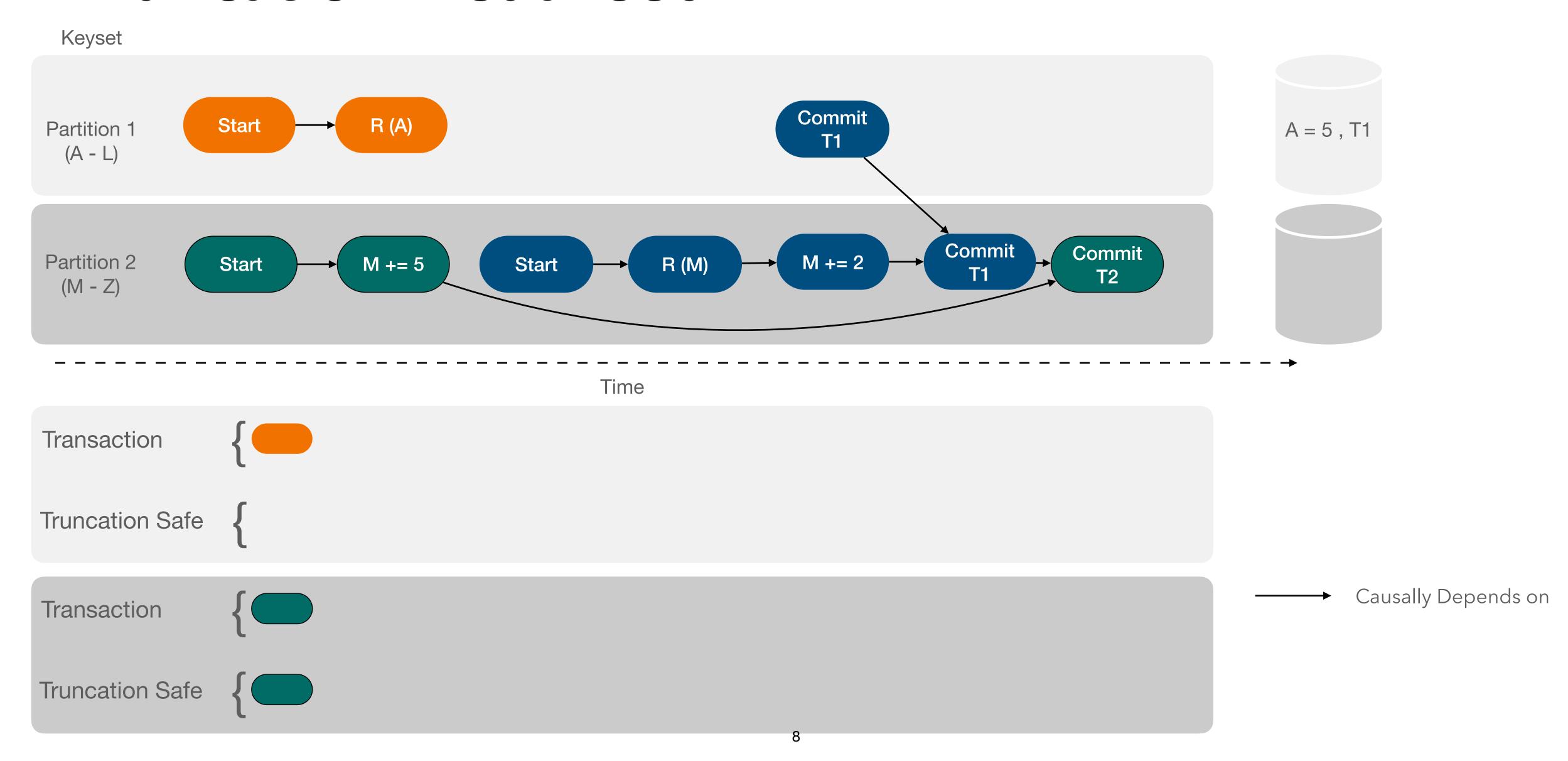


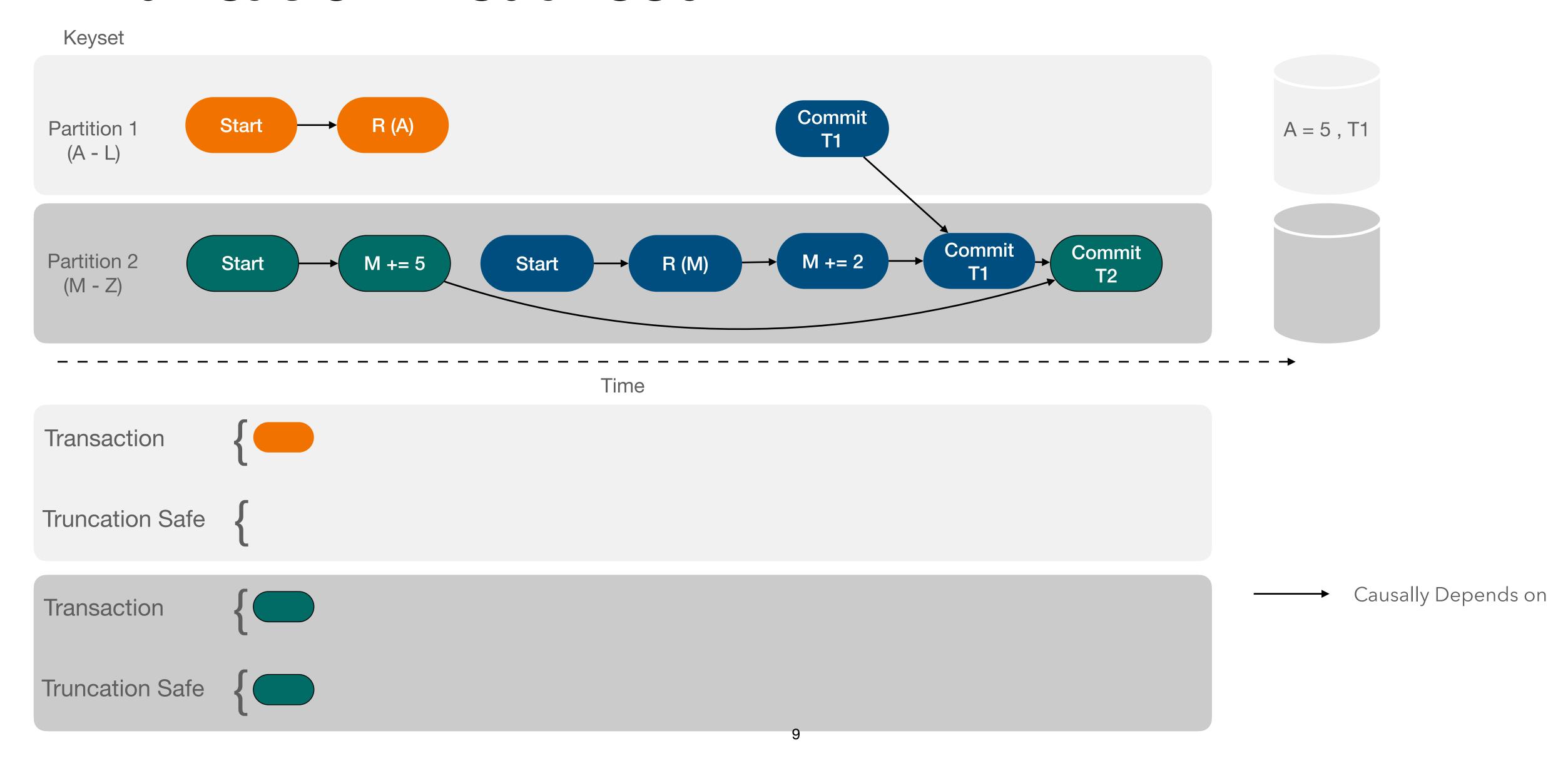


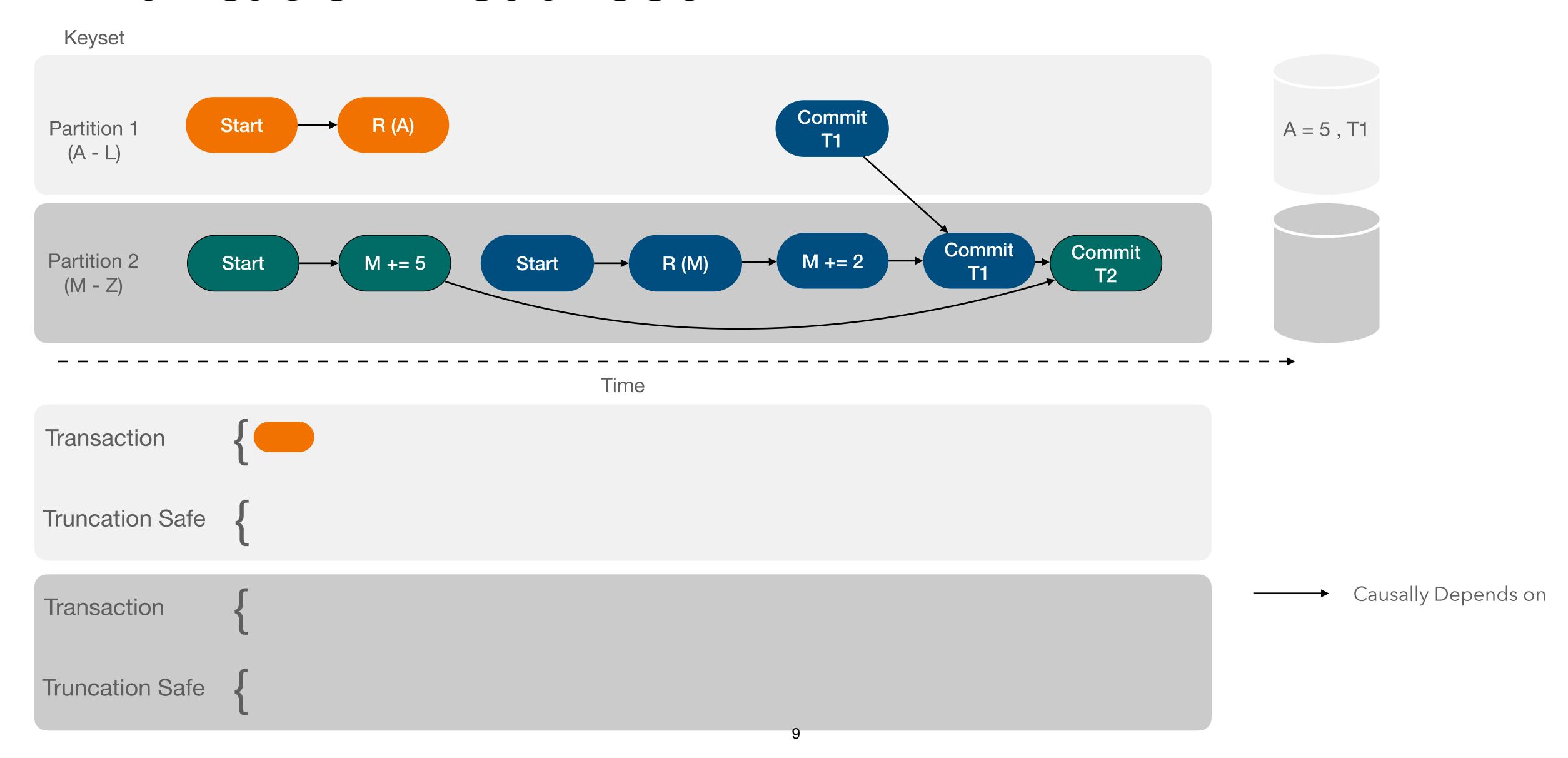


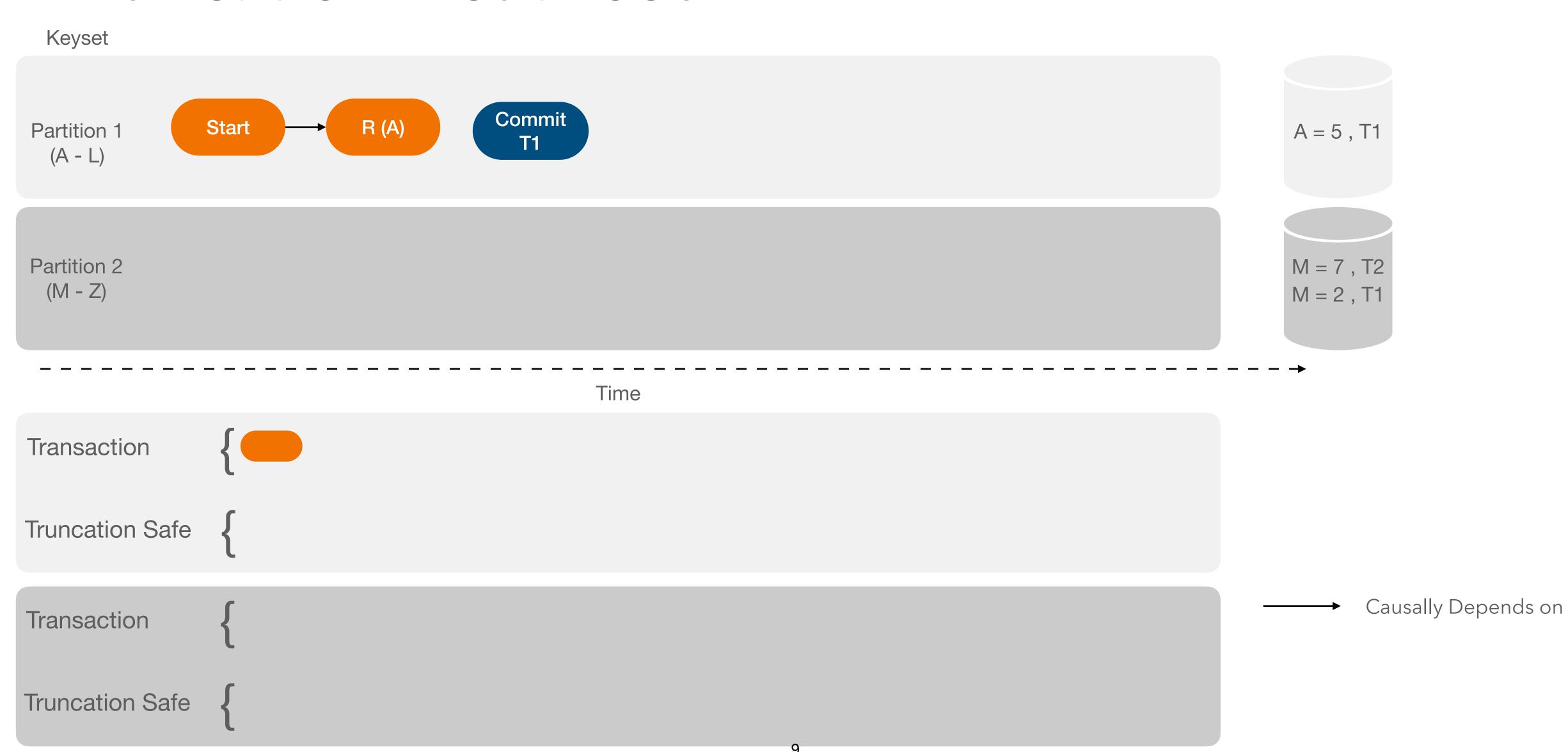


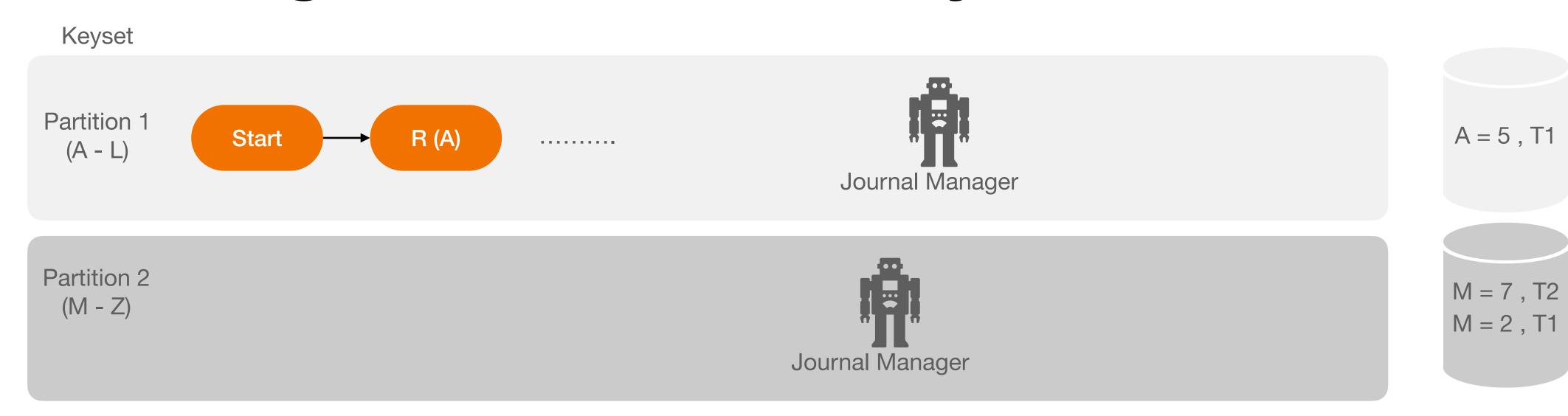




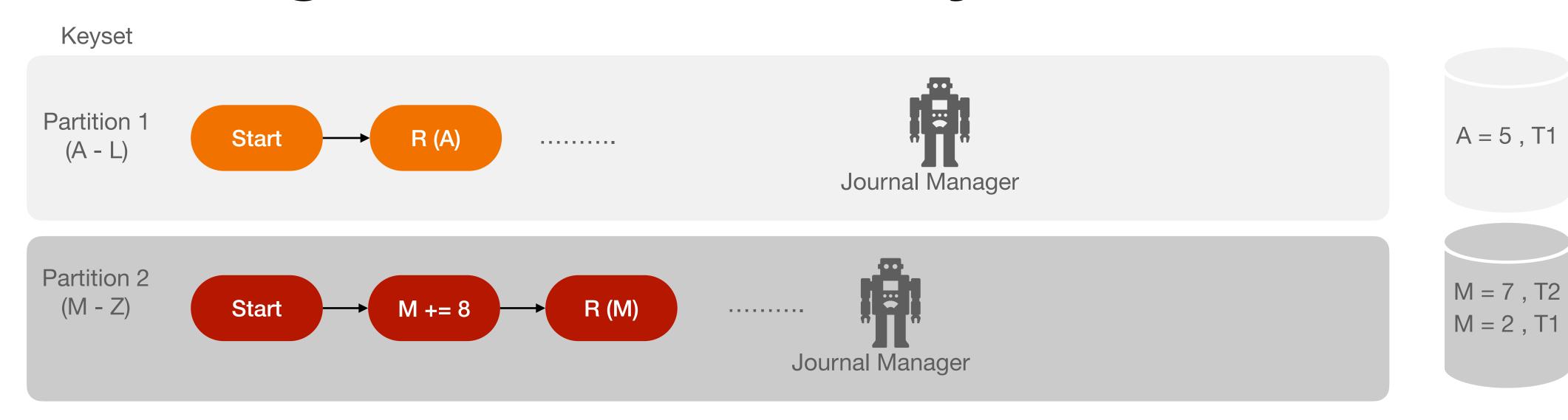




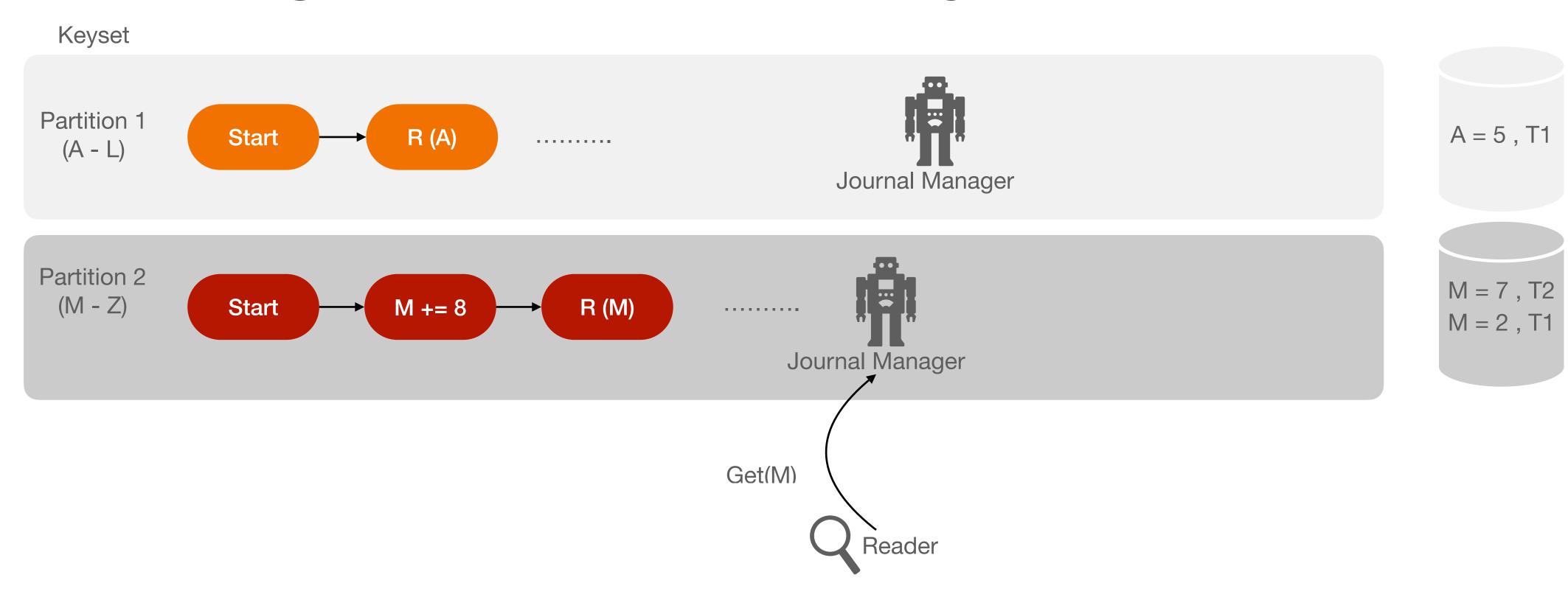


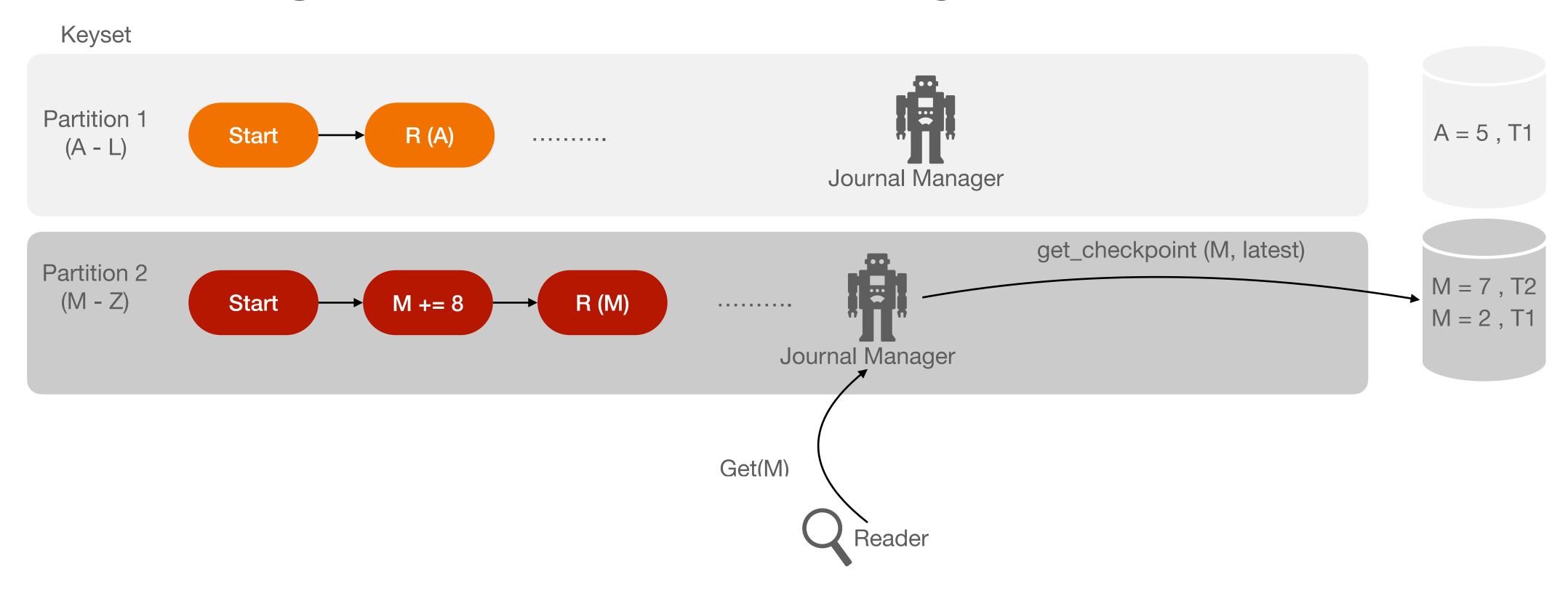


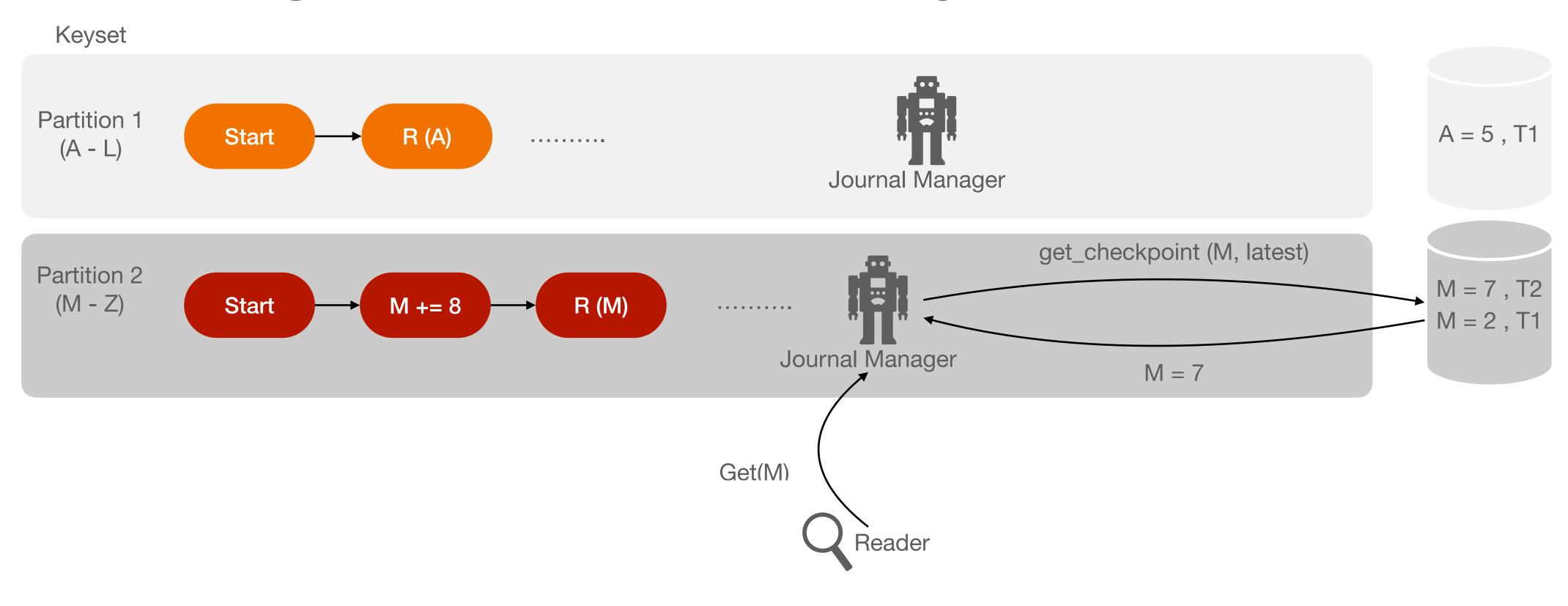


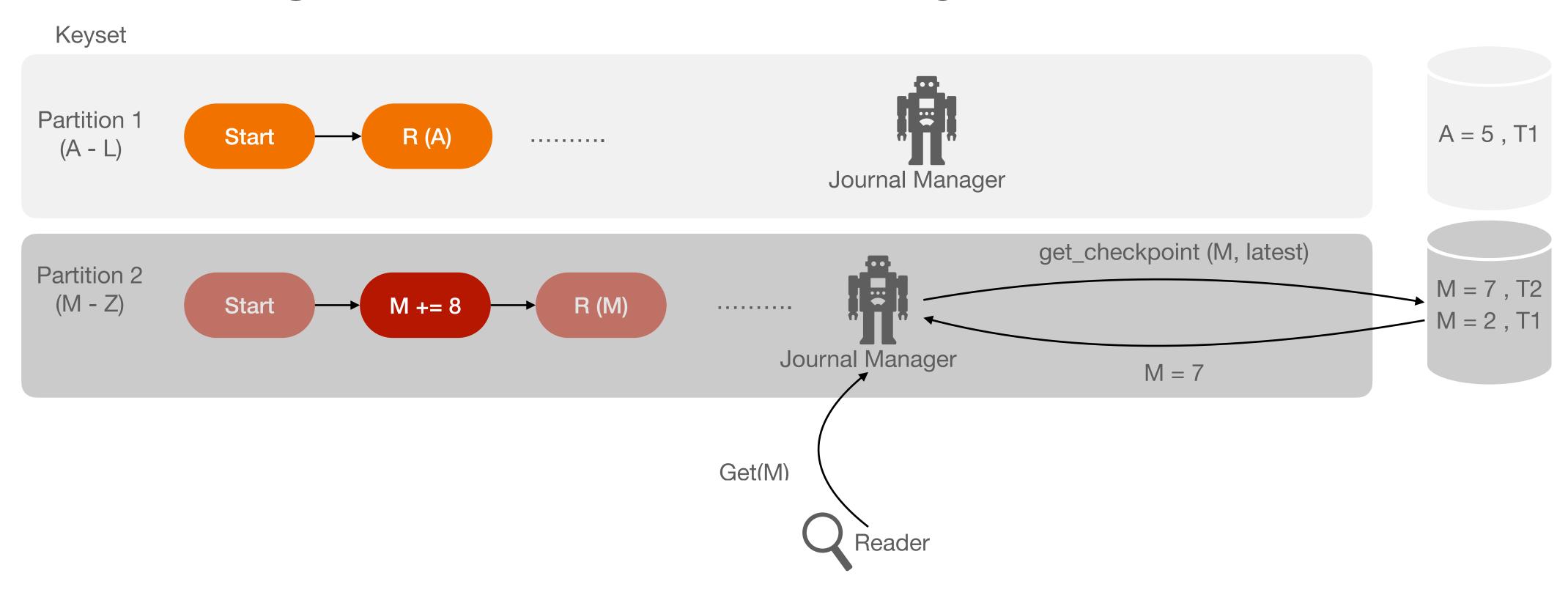


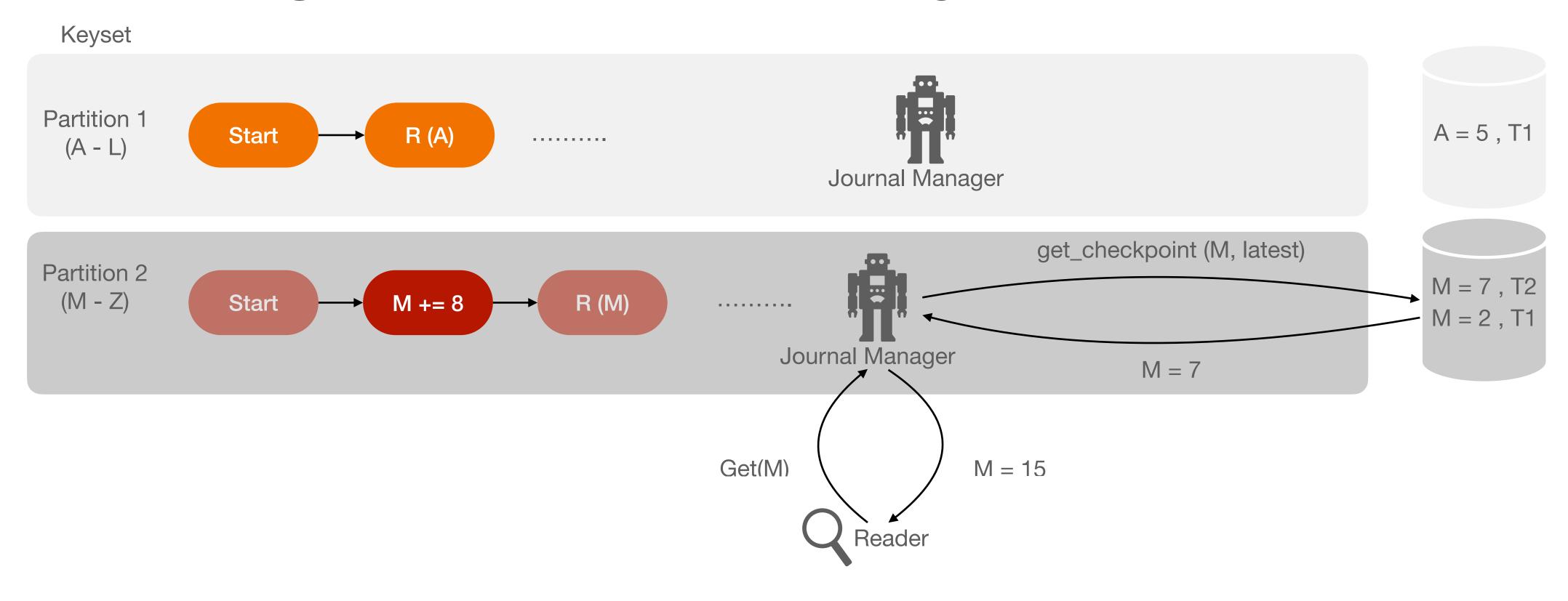








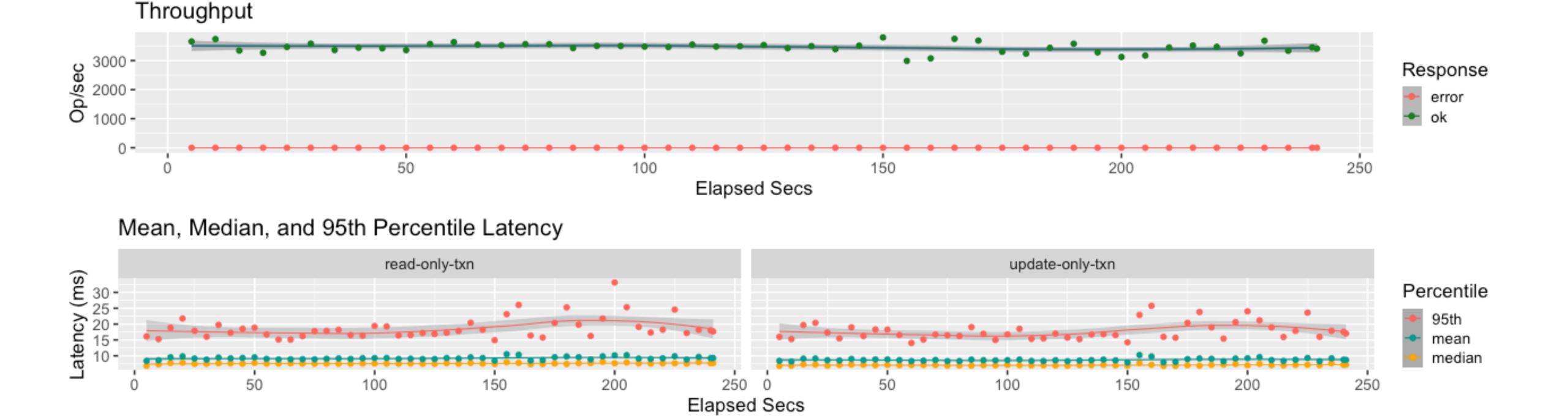




Implementation at a glance

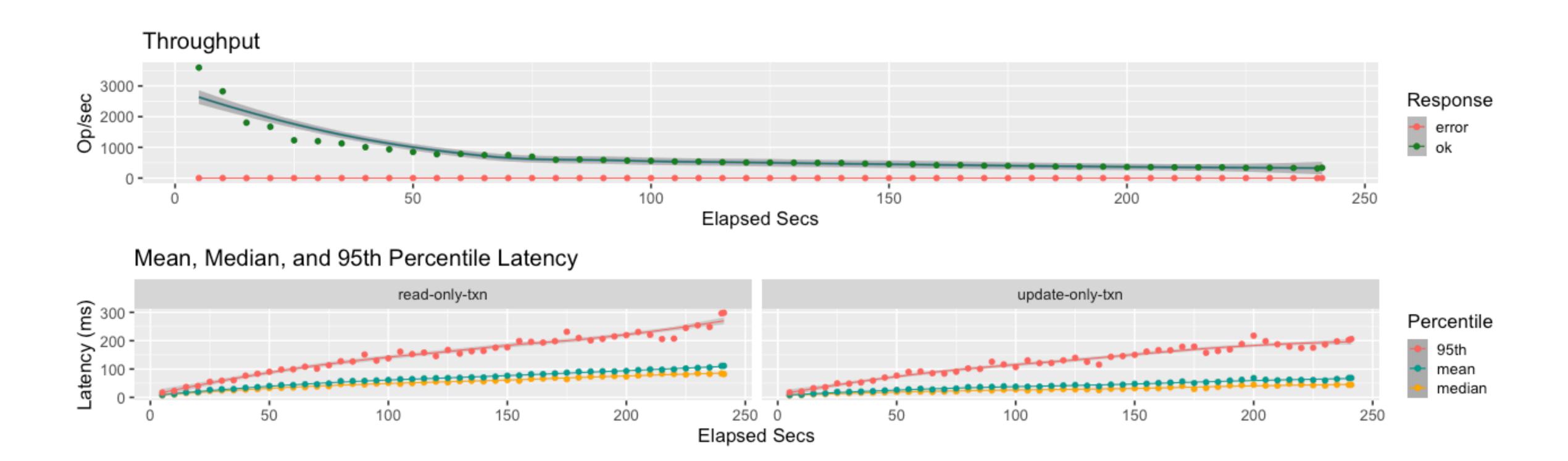
- Introducing a new library: Gingko.
- Restructuring the code into daemons and simplifying AntidoteDB:
 - Materialisation (Fill daemon)
 - Journal management (Gingko OP Log)
 - Journal Indexing (Index Daemon)
 - Checkpointing (Checkpoint Daemon)
 - Object Cache (Cache Daemon)
- Truncation is implemented via the Checkpoint Daemon.

Performance Benchmarks



Baseline Performance of the new Implementation

Performance Benchmarks



Observed decrease in throughput without Truncation

Conclusion

- In journal based data-stores, Journal grows with operational time.
- We need the journal to be bounded to manage persistent space requirement.
- Journal can be bounded by truncating the oldest records.
- Truncation needs to be done with transaction starts and commits in consideration.

Thank you!

New Architecture of a single DC in AntidoteDB

